

# A Unifying Formal Basis for the Sensoria Languages

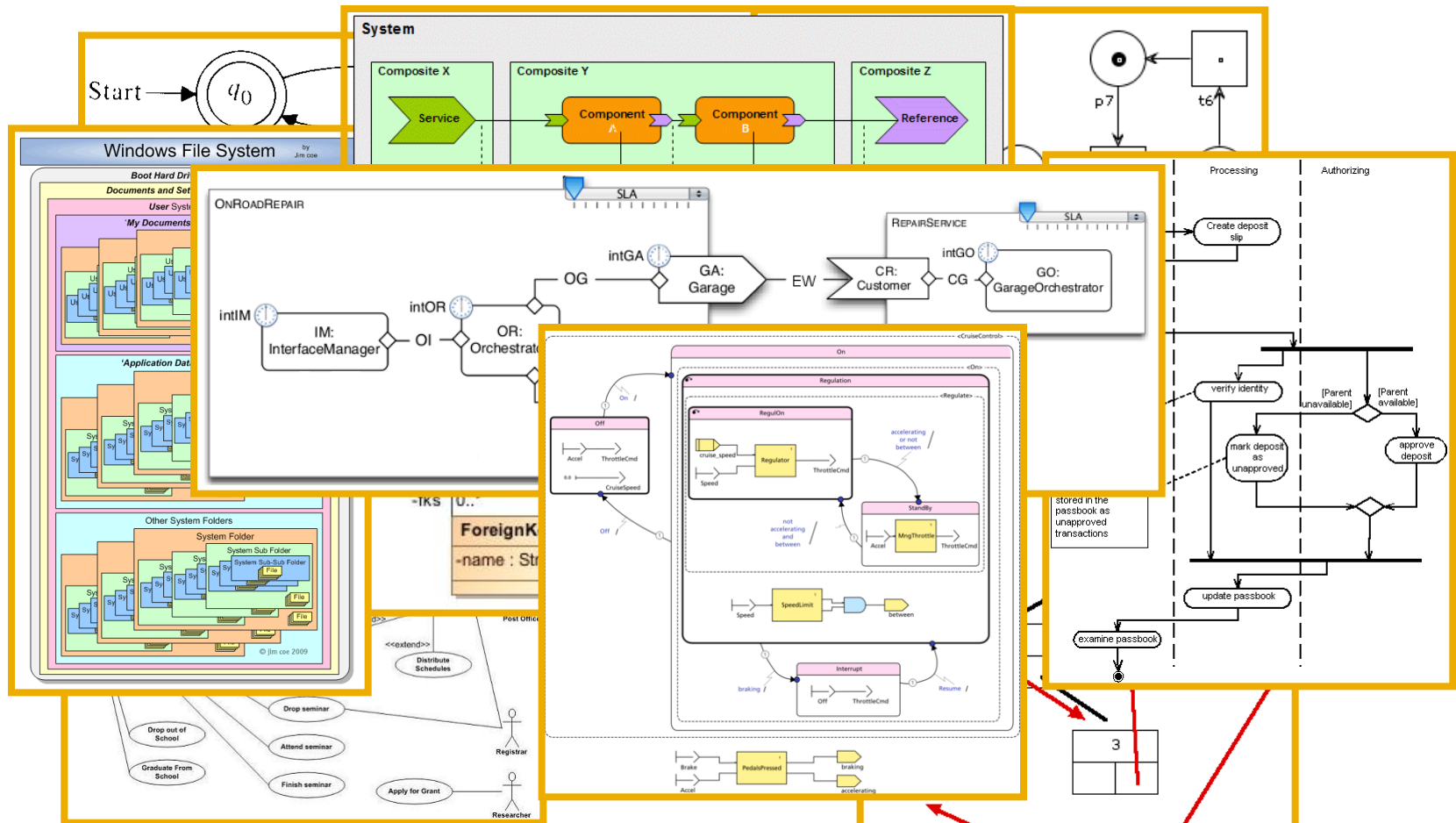
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- Introduction
- Modeling service oriented architectures
- Modeling service oriented calculi
- A general view

- SENSORIA Architecture level
  - High-level specification languages and frameworks
    - UML4SOA
    - SRML
    - Service Modes
  - Guaranteed reconfigurations
- SENSORIA Service-Oriented Calculi level
  - Process algebras & Process calculi approaches
    - SAGAS
    - CaSPiS
  - Sound and Complete graphical encoding

# Graphs Are Everywhere

- Use of diagrams / graphs is pervasive to Computer Science

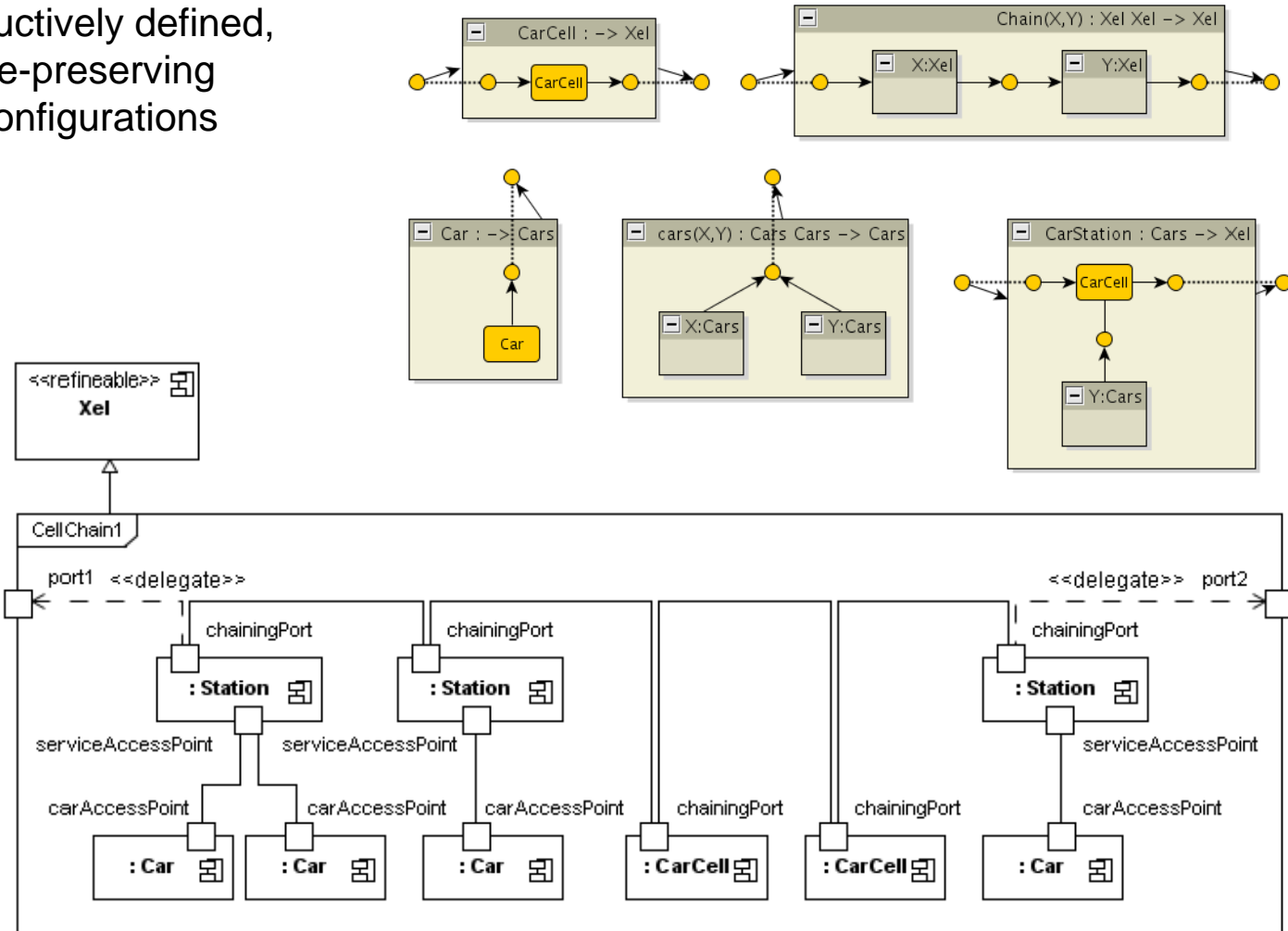


- ▶ ADR formulas:
  - ▶  $ADR = Designs + Term\ Rewriting$
  - ▶  $Designs = Typed\ Hierarchical\ Graphs\ (with\ Interfaces)$
- ▶ ADR ingredients:
  - ▶ **Sorts**: Vocabulary, Types (edge and node labels)
  - ▶ **Values**: Designs (hierarchical graphs with interfaces)
  - ▶ **Operations**: Graph-grammar-like rules
  - ▶ **Terms**: proofs of construction
  - ▶ **Terms (with variables)**: partial Designs, partial proofs
  - ▶ **Axioms**: properties of operations
  - ▶ **Membership predicates**: additional style rules
  - ▶ **Rewrite rules**: behaviour, reconfigurations
  - ▶ **Rewrite strategies**: style conformance, style analysis, etc.

- Why graph models?
  - More natural for distributed systems
  - Built-in structural axioms (e.g. name handling, AC axioms)
  - Uniform treatment of most ordinary process algebras (e.g. via the SHR approach)
- Why hierarchical?
  - Nested structures: ambients, block structure, sessions, transactions, etc.
  - Interaction between siblings, without referring to the closest common ancestor
- Yet another graph model?
  - Bigraphs by Robin Milner (2003)
    - place graph for localities and link graph for connectivity
    - Semantics via reduction rules and minimal contexts
  - Gs-monoidal graphs (gs-graphs) by Ferrari and Montanari (1997)
    - Based on gs-monoidal categories by Corradini and Gadducci
  - Top view & side view
- ADR metamodel and Design Algebra:
  - algebras of hierarchical graphs with graphical representations
- Applications to Sensoria
  - ADR modeling of UML4SOA, SRML, Modes
  - DA modeling of process calculi: transaction workflows (Sagas), service sessions (CaSPiS)
  - Reconcile Top view & Side view

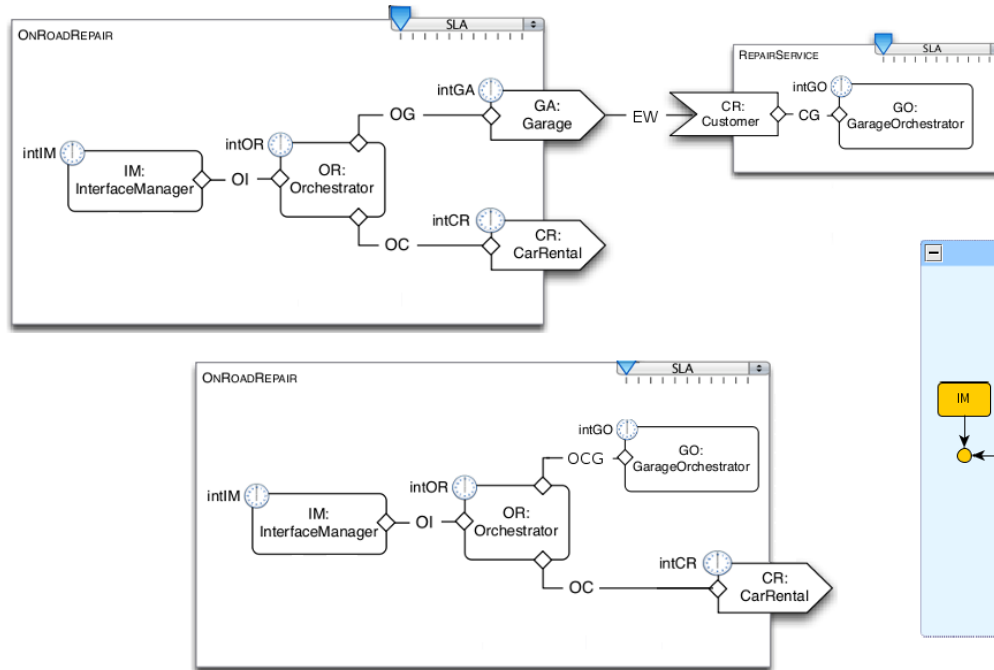
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Inductively defined,  
style-preserving  
reconfigurations

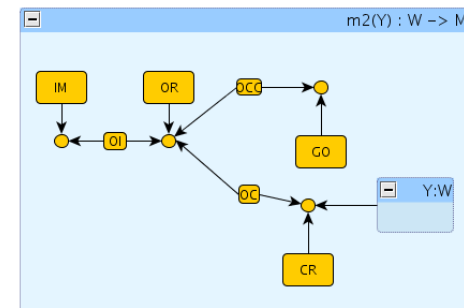
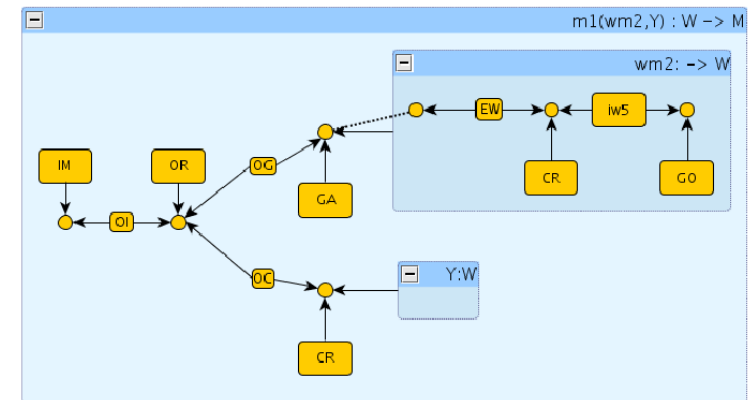




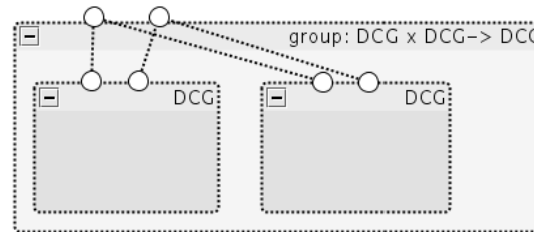
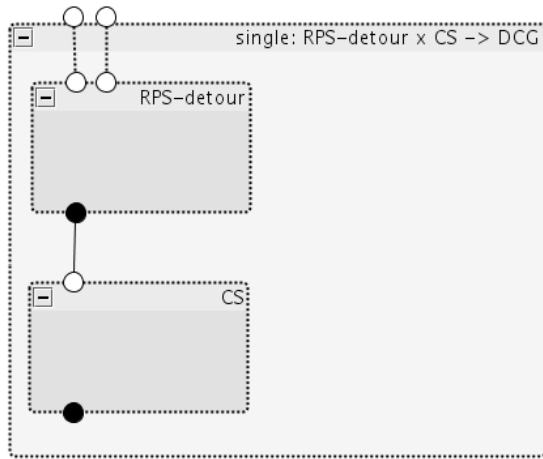
# ADR for SRML



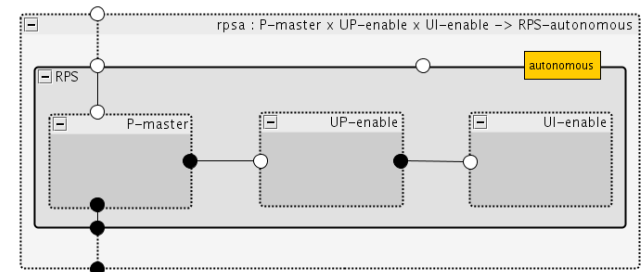
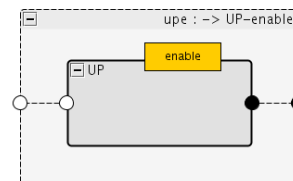
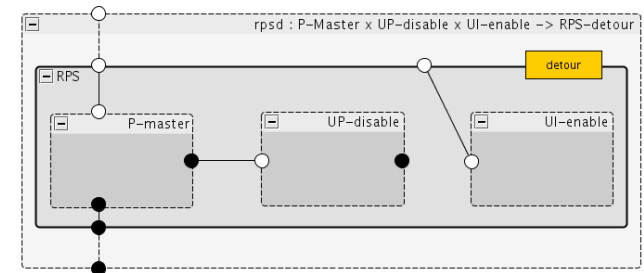
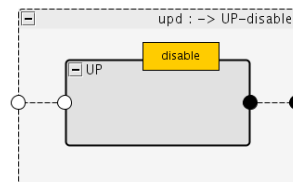
Inductively defined,  
well-formed,  
static & dynamic binding



Formal perational semantics via SHR rules  
(Synchronized Hyperedge Replacement)

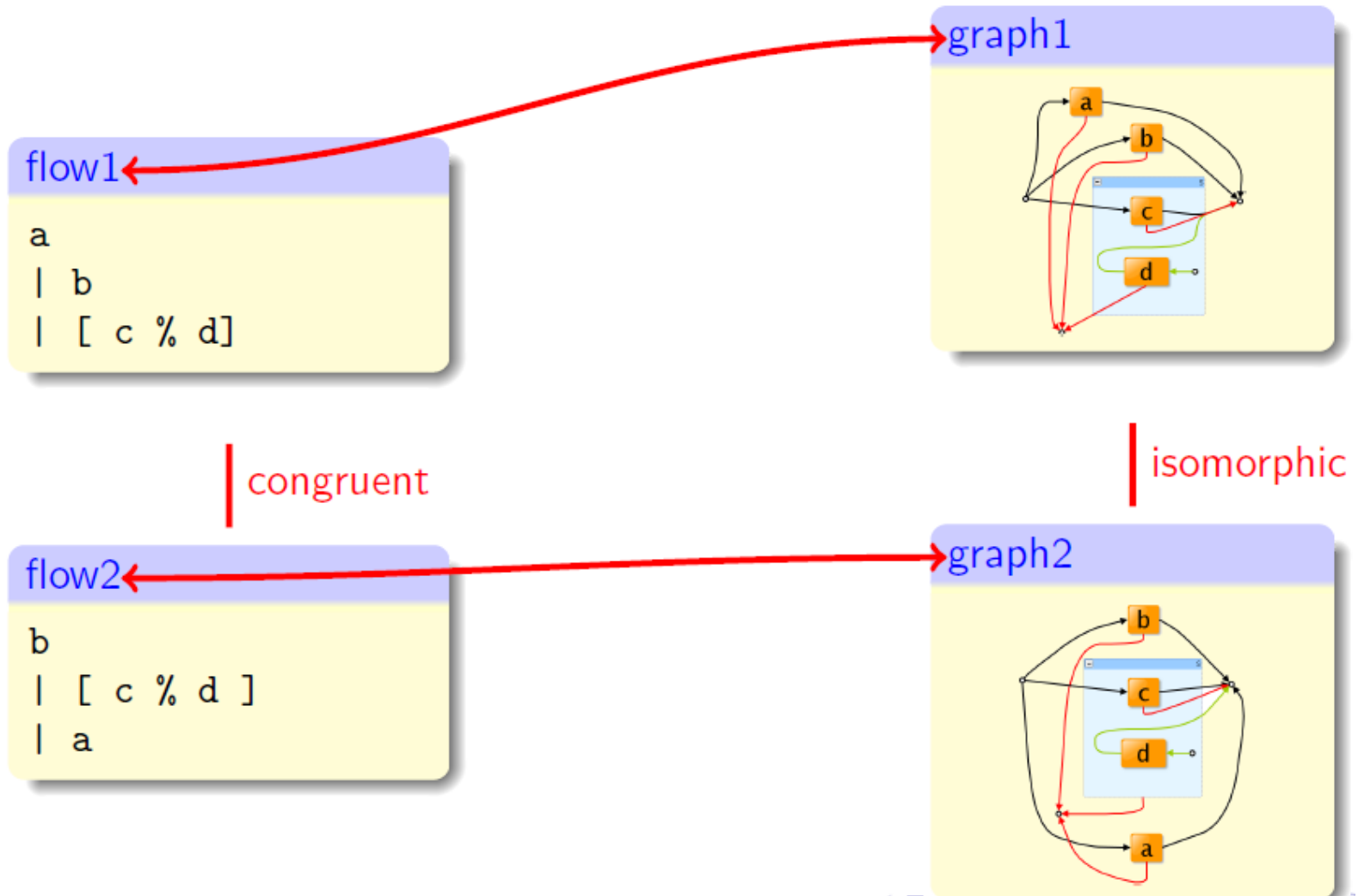


Inductively defined,  
mode-driven,  
style-preserving  
reconfigurations



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# The Goal: Sound and Complete Encoding Also for Hierarchical Graphs



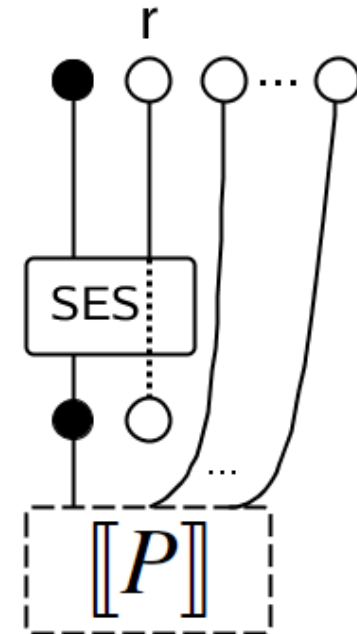
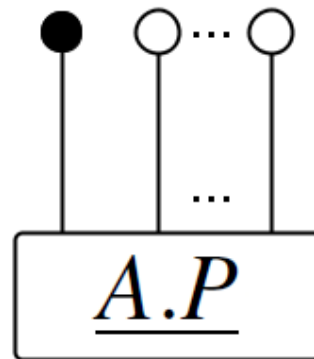
$$\begin{aligned} \mathbb{D} &::= L_{\bar{x}}[\mathbb{G}] \\ \mathbb{G} &::= \mathbf{0} \mid x \mid l(\bar{x}) \mid \mathbb{G} \mid \mathbb{G} \mid (\nu x)\mathbb{G} \mid \mathbb{D}(\bar{x}) \end{aligned}$$

$\mathbb{G} \mid \mathbb{H} \equiv \mathbb{H} \mid \mathbb{G}$
$\mathbb{G} \mid (\mathbb{H} \mid \mathbb{I}) \equiv (\mathbb{G} \mid \mathbb{H}) \mid \mathbb{I}$
$\mathbb{G} \mid \mathbf{0} \equiv \mathbb{G}$
$(\nu x)(\nu y)\mathbb{G} \equiv (\nu y)(\nu x)\mathbb{G}$
$(\nu x)\mathbf{0} \equiv \mathbf{0}$
$\mathbb{G} \mid (\nu x)\mathbb{H} \equiv (\nu x)(\mathbb{G} \mid \mathbb{H})$
$L_{\bar{x}}[\mathbb{G}] \equiv L_{\bar{y}}[\mathbb{G}\{y/\bar{x}\}]$
$(\nu x)\mathbb{G} \equiv (\nu y)\mathbb{G}\{y/x\}$
$x \mid \mathbb{G} \equiv \mathbb{G}$
$L_{\bar{x}}[z \mid \mathbb{G}](\bar{y}) \equiv z \mid L_{\bar{x}}[\mathbb{G}](\bar{y})$

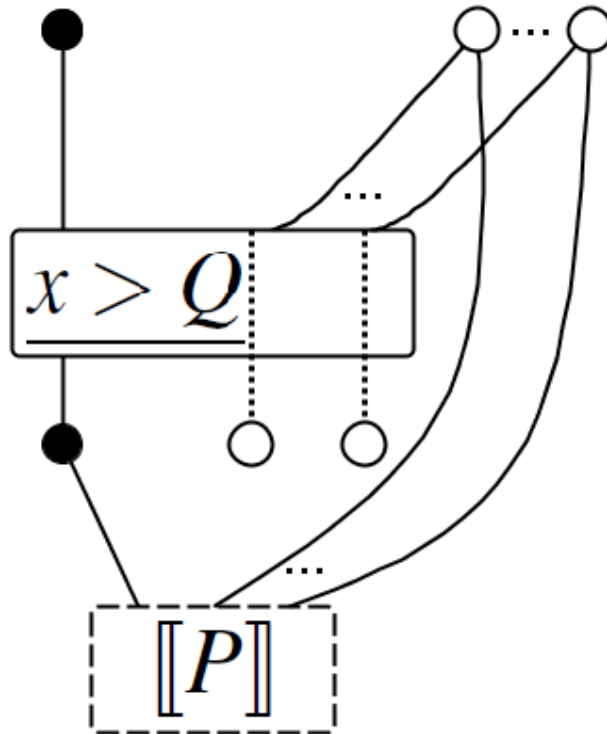
(DA1)	$fn(\mathbf{0}) = \emptyset$
(DA2)	$fn(l(\bar{x})) = [\bar{x}]$
(DA3)	$fn((\nu x)\mathbb{G}) = fn(\mathbb{G}) \setminus \{x\}$
(DA4)	$fn(L_{\bar{x}}[\mathbb{G}]) = fn(\mathbb{G}) \setminus [\bar{x}]$
(DA5)	$fn(x) = x$
(DA6)	$fn(\mathbb{G} \mid \mathbb{H}) = fn(\mathbb{G}) \cup fn(\mathbb{H})$
(DA7)	$fn(\mathbb{D}(\bar{x})) = fn(\mathbb{D}) \cup [\bar{x}]$
(DA8)	$if\ x \notin fn(\mathbb{G})$
(DA9)	$if\ [y] \cap fn(\mathbb{G}) = \emptyset$
(DA10)	$if\ y \notin fn(\mathbb{G})$
	$if\ x \in fn(\mathbb{G})$
	$if\ z \notin [\bar{x}]$

Theorem: Sound and Complete axioms for hierarchical graphs  
(also surjective map from terms to hierarchical graphs)

Theorem: Sound and Complete axioms for gs-graphs

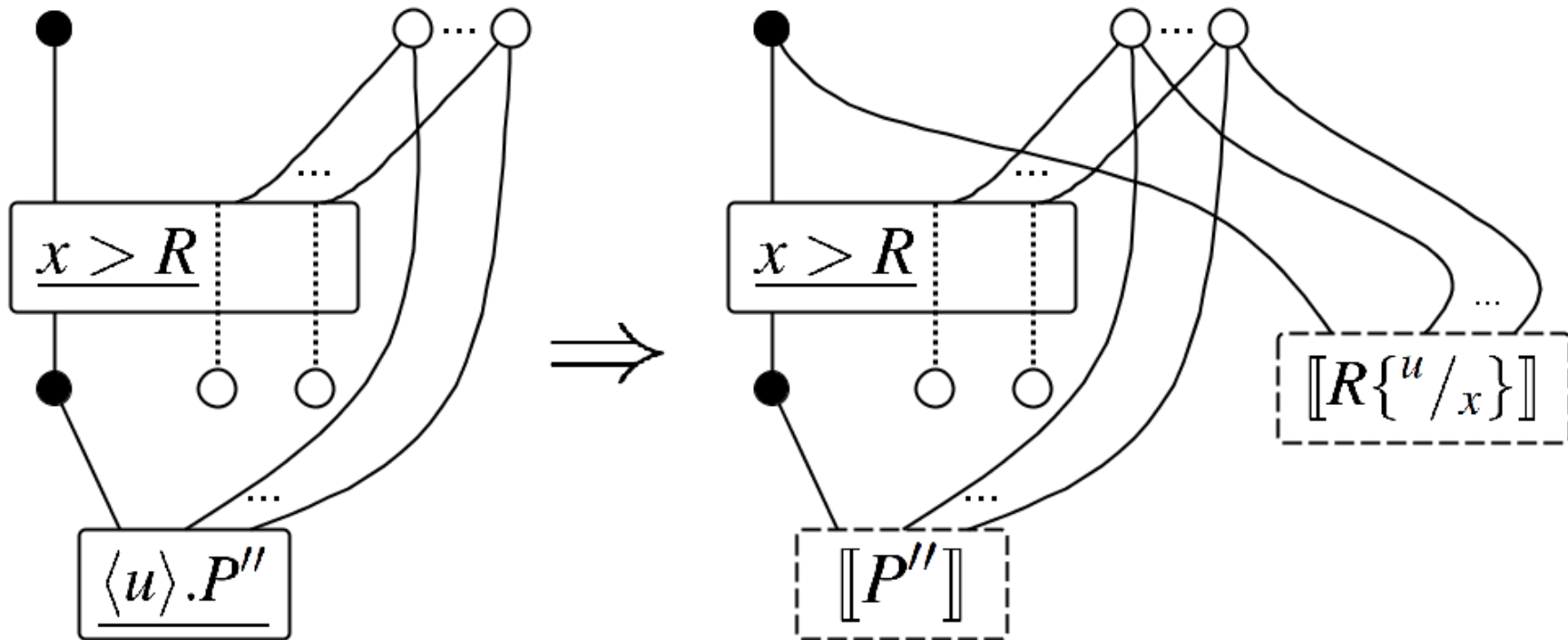


$$\begin{aligned} \llbracket A.P \rrbracket &\stackrel{\text{def}}{=} \underline{A.P} \langle \llbracket fn(A.P) \rrbracket \rangle \\ \llbracket r \triangleright P \rrbracket &\stackrel{\text{def}}{=} SES[ \llbracket P \rrbracket ] \langle r \rangle \end{aligned}$$



pipeline construct

$$\llbracket P > (?x)Q \rrbracket \stackrel{\text{def}}{=} \underline{x > Q} [ \llbracket P \rrbracket ] \langle [fn(Q) \setminus \{x\}] \rangle$$



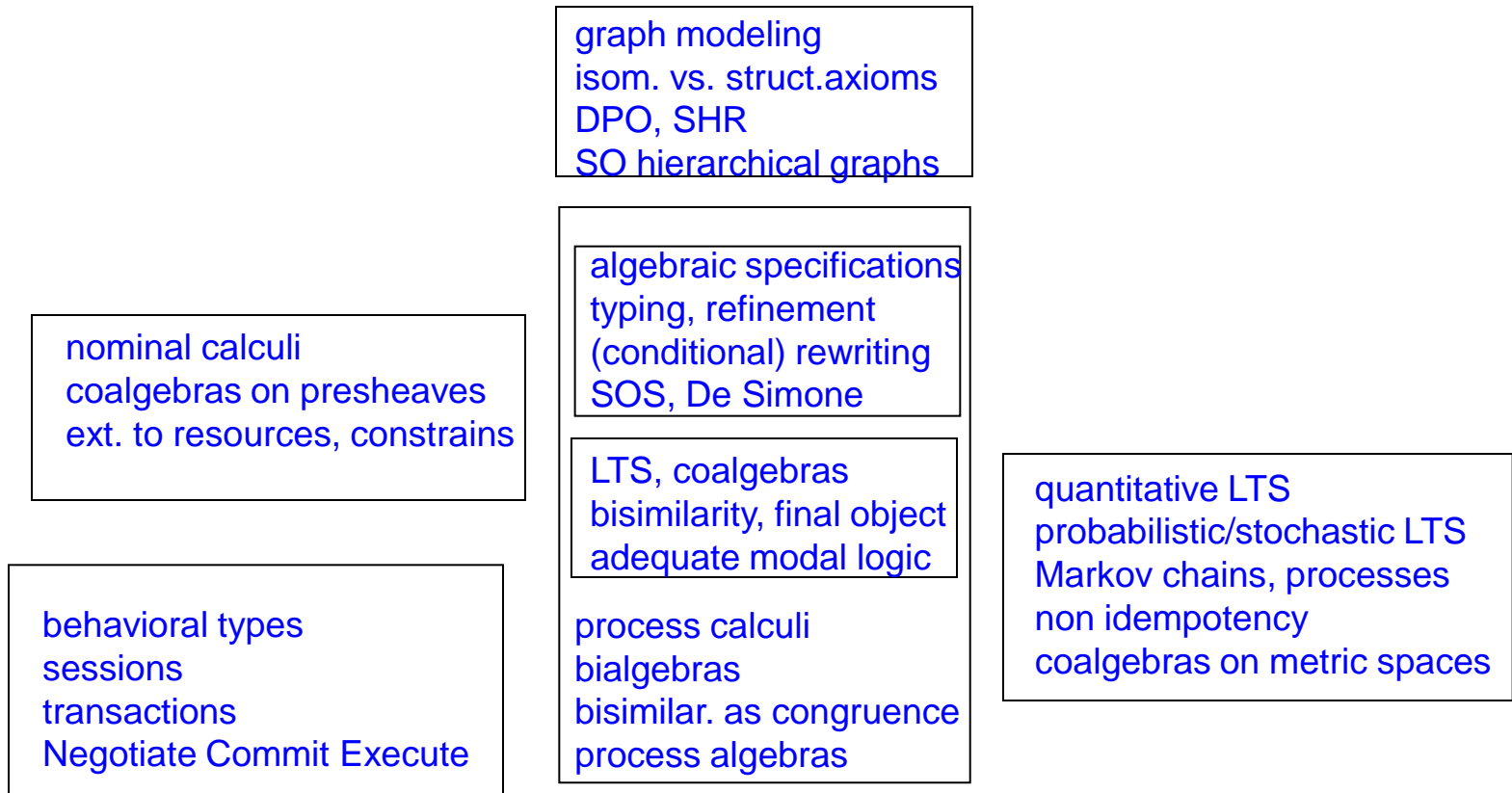
$$P \equiv C[ (P' | \langle u \rangle . P'') > (?x).R ]$$

$$Q \equiv C[ R\{u/x\} \mid ((P' | P'') > (?x).R) ]$$

(PipelineSync)



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# THANKS FOR THE ATTENTION

